### The Functorial Difference Operator

Robert Paré (Dalhousie University) pare@mathstat.dal.ca

Topos Institute Colloquium

January 11, 2024

### What is an endofunctor of Set like?

• Polynomials  $F(X) = \sum_{i \in I} X^{A_i}$  (Kock 2009, Spivak *et al.* – see [3, 5])

• Analytic functors 
$$F(X) = \int_{0}^{n} X^{n} \times C_{n}$$
 (Joyal 1981 [2])

Monads (Manes 2002 [4], Szawiel Zawadowski 2015 [7])

• Reduced powers  $F(X) = X^{\mathscr{F}}$ ,  $\mathscr{F}$  filter (Blass 1976 [1])

### The structure of endofunctors of Set

• We study  $F: \mathbf{Set} \longrightarrow \mathbf{Set}$  by perturbing X and measuring the change in F(X).

Example  $FX = X^3$ An element of  $F(X + 1) = (X + \{*\})^3$ 

• Going from X to X+1, F gains  $3X^2+3X+1$  elements.

### Tautness

# Definition (Manes 2002 [4])

A functor is *taut* if it preserves inverse images

A natural transformation  $t: F \longrightarrow G$  is *taut* if the naturality squares for monomorphisms are pullbacks

$$\begin{array}{ccc} FA_{0} \longrightarrow FA \\ A_{0} \rightarrowtail A & \Longrightarrow & tA_{0} \bigvee \begin{tabular}{c} \mathsf{Pb} \\ \mathsf{G}A_{0} & \longrightarrow \begin{tabular}{c} \mathsf{G}A_{0} \\ \mathsf{G}A_{0} & \longrightarrow \begin{tabular}{c} \mathsf{G$$

# The plenitude of tautness

There are plenty of taut functors:

- Polynomial functors
- Analytic functors
- Reduced powers
- Left exact functors
- Functors Set —> Set that preserve binary coproducts

They are closed under a variety of operations.

We get a sub-2-category  $\mathcal{T}aut$  of  $\mathcal{C}at$ , whose objects are categories with inverse images, 1-cells are taut functors, 2-cells taut natural transformations.

# Limits

### Proposition

Assume that **B** has I-limits and let  $\Gamma: I \longrightarrow Cat(\mathbf{A}, \mathbf{B})$  be a diagram.

- (1) If  $\Gamma I$  is taut for all I, then  $\lim_{I \to I} \Gamma I$  is taut.
- (2) If, furthermore, I is non-empty and connected and Γ(i) is a taut transformation for all i: I→I' then lim<sub>I</sub> Γ(I) is the limit in *Taut*(A,B), i.e., the inclusion *Taut*(A,B) → Cat(A,B) creates connected limits.

### Example

If  $F, G: \mathbf{A} \longrightarrow \mathbf{B}$  are taut and  $\mathbf{B}$  has finite products, then  $F \times G$  is taut, but the projections are not.

The constant functor 1 is taut but not terminal in  $\mathcal{T}aut(\mathbf{A}, \mathbf{B})$ .

### Confluence

Definition

I is *confluent* if every span in I can be completed to a commutative square



#### Theorem

I-colimits commute with inverse images in Set if and only if I is confluent.

Remark This means  $\lim$ : Set<sup>I</sup>  $\rightarrow$  Set is taut.

# Colimits

### Theorem

- (1) Confluent colimits of taut functors in Cat(A, Set) are taut.
- (2) The inclusion  $\mathcal{T}aut(\mathbf{A}, \mathbf{Set}) \longrightarrow \mathcal{C}at(\mathbf{A}, \mathbf{Set})$  creates confluent colimits.

### Example

Coproducts, filtered colimits, quotients by a group action, are all confluent.

# Polynomials

$$P(X) = \sum_{i \in I} X^{A_i} \text{ is taut.}$$

A morphism of polynomials is a natural transformation  $t(X): P(X) \longrightarrow Q(X)$ .

If 
$$Q(X) = \sum_{j \in J} X^{B_j}$$
, morphisms  $P(X) \longrightarrow Q(X)$  correspond to

$$\alpha: I \longrightarrow J, \quad \langle f_i: B_{\alpha(i)} \longrightarrow A_i \rangle_i$$

f is vertical if  $\alpha$  is an identity, cartesian if all the  $f_i$  are isomorphisms.

#### Proposition

t is taut if and only if all the  $f_i$  are epimorphisms.

# Analytic functors

- Species F: Bij --> Set
- Analytic functor



(Left Kan extension)

$$\widetilde{F}(X) = \int^{n \in \mathbb{N}} X^n \times F(n)$$
  
$$\cong \lim_{a \in F(n)} X^n$$
  
$$\cong \sum_{n \in \mathbb{N}} (X^n \times F(n)) / S_n$$

Proposition  $\tilde{F}$  is taut.

### Reduced powers

- Filter  $\mathscr{F} \subseteq 2^A$  closed under finite intersections upclosed
- Reduced power  $X^{\mathscr{F}}$

$$\begin{split} X^A / &\sim \qquad (f \sim g \Leftrightarrow \{a \in A \mid f(a) = g(a)\} \in \mathscr{F}) \\ &\cong \lim_{B \in \mathscr{F}} X^B \end{split}$$

Proposition  $X^{\mathscr{F}}$  is taut.

**Note:**  $X^{\mathscr{F}}$  is not an analytic functor, unless  $\mathscr{F}$  is principal  $(X^{\langle A_0 \rangle} \cong X^{A_0})$ .

# Monads

- The free monoid monad  $1 + X + X^2 + \dots$  is taut.
- The free commutative monoid monad  $1 + X + X^2/S_2 + X^3/S_3 + ...$  is taut.
- Manes (2002 [4]) Collection monads are finitary taut monads.
- The free abelian group monad is not taut.

• Płonka (1967) [6] - Balanced equations (same variables on both sides).

• Szawiel/Zawadowski (2015) [7] - A finitary monad is taut if and only if it can be defined by balanced equations.

### The difference operator

For  $F: \mathbf{Set} \longrightarrow \mathbf{Set}$  define  $\Delta[F]: \mathbf{Set} \longrightarrow \mathbf{Set}$  by  $\Delta[F](X) = F(X+1) \setminus F(X)$ .

### Example

 $\Delta[C] = 0 \quad \Delta[X] = 1$ 

Proposition

If F is taut, then  $\Delta[F](X)$  is a taut subfunctor of F(X+1).

Everything hinges on the following fact:

For a diagram of sets and functions

$$\begin{array}{c|c} A_{0} \longrightarrow A \\ f_{0} \downarrow & (*) & \downarrow f \\ B_{0} \longrightarrow B \end{array}$$

f restricts to the complements  $A'_0$  and  $B'_0$  iff (\*) is a pullback.

# Colimits

Let **Taut** =  $\mathcal{T}aut$ (**Set**, **Set**).

Proposition

 $\Delta$  is a functor, the difference operator,

 $\Delta$ : Taut  $\rightarrow$  Taut.

It preserves confluent colimits

 $\Delta[\varinjlim_I \Gamma I] \cong \varinjlim_I \Delta[\Gamma I] \ .$ 

Corollary (1)  $\Delta[CF] \cong C\Delta[F]$ . (2)  $\Delta[F+G] \cong \Delta[F] + \Delta[G]$ .

# Limits

 $\begin{array}{l} \mathsf{Proposition} \\ \Delta[F \times G] \cong (\Delta[F] \times G) + (F \times \Delta[G]) + (\Delta[F] \times \Delta[G]). \end{array}$ 

More generally:

Proposition

$$\Delta[\prod_{i\in I} F_i] \cong \sum_{J\subsetneq I} (\prod_{j\in J} F_j) \times (\prod_{k\notin J} \Delta[F_k]).$$

### Theorem

 $\Delta$  preserves non-empty connected limits:

$$\Delta[\varprojlim_I \Gamma I] \cong \varprojlim_I \Delta[\Gamma I].$$

Polynomials

•  $\Delta[X^A] \cong \sum_{B \subsetneq A} X^B$ 

### Proposition

If P(X) is a polynomial functor, then so is  $\Delta[P(X)]$ . For  $P(X) = \sum_{i \in I} X^{A_i}$ ,

$$\Delta[P(x)] = \sum_{j \in J} X^{B_j}$$

where  $J = \{(i, B) \mid i \in I, B \subsetneq A_i\}$  and for  $j = (i, B), \quad B_j = B.$ 

• 
$$\Delta[X^n] \cong \sum_{k=0}^{n-1} \binom{n}{k} X^k$$

Proposition

If F(X) is a power series functor  $\sum_{n=0}^{\infty} C_n X^n$ , then  $\Delta[F(X)]$  is also a power series  $\sum_{n=0}^{\infty} D_n X^n$  where  $D_n = \sum_{k=1}^{\infty} {n+k \choose k} C_k$ .

# Analytic functors

### Proposition

If  $\tilde{F}(X)$  is an analytic functor corresponding to the species  $F: \operatorname{Bij} \longrightarrow \operatorname{Set}$ , then  $\Delta[\tilde{F}(X)]$  is also analytic, corresponding to the species

$$G(n) = \int^{k \in \mathbb{N}^+} F(n+k).$$

• A *G*-structure of cardinality n consists of a positive integer k and an equivalence class of *F*-structures of cardinality n + k. Two such structures are equivalent if one is transformed into the other by a bijection fixing the first n elements.

### Reduced powers

A filter  $\mathcal{F}$  on A induces an equivalence relation on subsets of A

$$\begin{array}{ll} B \sim C & \Leftrightarrow & \{a \in A \mid a \in B \Leftrightarrow a \in C\} \in \mathscr{F} \\ & \Leftrightarrow & (B \cap C) \cup (B' \cap C') \in \mathscr{F}. \end{array}$$

For every  $B \subseteq A$ , let  $\mathscr{F}_B = \{C \subseteq B \mid C \cup B' \in \mathscr{F}\}.$ 

#### Proposition

$$\mathscr{F}_B$$
 is a filter on  $B$  and  $\Delta[X^{\mathscr{F}}] \cong \sum_{[B] \neq [A]} X^{\mathscr{F}_B}.$ 

(The sum is over all equivalence classes not equal to [A], one summand for each class.)

### Lax chain rule

Theorem

For taut functors F and G there is a taut natural transformation

 $\gamma_{G,F} \colon (\Delta[G] \circ F) \times \Delta[F] \longrightarrow \Delta[G \circ F]$ 



# Newton series

• For  $f: \mathbb{R} \longrightarrow \mathbb{R}$  its *Newton series* is

$$\sum_{n=0}^{\infty} \frac{\Delta^n[f](0)}{n!} x^{\downarrow n} = \sum_{n=0}^{\infty} \Delta^n[f](0) \binom{x}{n}.$$

$$-x^{\downarrow n} = \text{falling power } x(x-1)\cdots(x-n+1)$$

$$-\binom{x}{n} =$$
 "binomial coefficient"  $\frac{x(x-1)\cdots(x-n+1)}{n!}$ 

 $-\Delta^n[f]$  is iterated difference

$$\begin{split} &\Delta^0[f](x) = f(x) \\ &\Delta^1[f](x) = f(x+1) - f(x) \\ &\Delta^2[f](x) = f(x+2) - 2f(x+1) + f(x) \\ &\Delta^3[f](x) = f(x+3) - 3f(x+2) + 3f(x+1) - f(x) \\ & etc. \end{split}$$

### Iterated difference

Proposition

 $\Delta^{n}[F](X) = \left\{ a \in F(X+n) \mid a \notin F(X+k) \text{ for any proper subset } k \subsetneq n \right\}.$ 

 $S_n$  acts on  $\Delta^n[F](X)$  giving a species  $\Delta^*[F](0)$  and a corresponding analytic functor

$$\sum_{n=0}^{\infty} \left( X^n \times \Delta^n[F](0) \right) / S_n.$$

But this won't give F(X), even for polynomials. However  $\Delta^n[F](X)$  has more "symmetries".

#### Proposition

If  $e: n \longrightarrow m$  is onto,  $F(X + e): F(X + n) \longrightarrow F(X + m)$  restricts to

 $\Delta^{e}[F](X) \colon \Delta^{n}[F](X) \longrightarrow \Delta^{m}[F](X).$ 

### Soft species

### Definition

Let **Surj** be the category of finite cardinals and surjections. A *soft species* is a functor  $C: Surj \rightarrow Set$ . It determines a *soft analytic functor* (semi-analytic in [7]) by left Kan extension along the inclusion of **Surj** into **Set**:



### Proposition

Analytic functors are soft analytic. Soft analytic functors are taut.

### Soft analytic functors

For C: Surj  $\rightarrow$  Set, an element of  $\widetilde{C}(X)$  is an equivalence class

 $[a \in C(n), f: n \longrightarrow X].$ 

Factor *f*:



So every equivalence class has a representation with f monic.

$$\widetilde{C}(X) \cong \sum_{n=0}^{\infty} \left( C(n) \times \mathsf{Mono}(n, X) \right) / S_n \cong \sum_{n=0}^{\infty} C(n) \times \binom{X}{n}$$

but only as sets!

### Newton series

For  $F: \mathbf{Set} \longrightarrow \mathbf{Set}$  a taut functor, the sets  $\Delta^n[F](0)$  extend to a soft species

 $\Delta^*[F](0):$  Surj $\longrightarrow$ Set.

The corresponding soft analytic functor

$$\overline{F}(X) = \widetilde{\Delta^*[F](0)} = \int^{n \in \operatorname{Surj}} \Delta^n[F](0) \times X^n$$

is the *Newton series* of *F*.

As sets

$$\overline{F}(X) \cong \sum_{n=0}^{\infty} \left( \Delta^n[F](0) \times \mathsf{Mono}(n, X) \right) / S_n \cong \sum_{n=0}^{\infty} \Delta^n[F](0) \times \binom{X}{n}.$$

Compare with:

$$\sum_{n=0}^{\infty} \frac{\Delta^n[f](0)}{n!} x^{\downarrow n} = \sum_{n=0}^{\infty} \Delta^n[f](0) \binom{x}{n}.$$

# Fundamental theorem of functorial differences

Let SoftSp be the category Set<sup>Surj</sup> of soft species and natural transformations.

Theorem (1)  $\widetilde{()}$  gives a functor SoftSp $\rightarrow$ Taut.

(2)  $F \mapsto \langle \Delta^n[F](0) \rangle_n$  gives a functor  $\Delta^*$ : Taut  $\longrightarrow$  SoftSp.

(3)  $\widetilde{()}$  is left adjoint to  $\Delta^*$ .

(4) The unit is an isomorphism  $C \xrightarrow{\cong} \Delta^*[\widetilde{C}](0)$ .

### Corollary

The Newton sum of a soft analytic functor "converges to it".

## Conclusion

We've:

- Identified taut functors as the context to develop a functorial calculus of differences.
- Discovered confluent colimits which are central.
- Generalized the sum and product rules to colimits and limits.
- Established a lax chain rule.
- Expressed Newton summation as a left adjoint.

A multivariable version is in preparation.

Thank you!

# References

[1] Andreas Blass.

Exact functors and measurable cardinals. *Pacific J. Math.*, 63(2):335–346, 1976.

[2] André Joyal.

Une théorie combinatoire des séries formelles. Advances in Mathematics, 42(1):1–82, October 1981.

[3] Joachim Kock.

#### Notes on polynomial functors.

http://mat.uab.cat/~kock/cat/polynomial.html, August 16 2016.

[4] Ernest G. Manes.

Taut monads and *T*0-spaces. *Theoretical Computer Science*, 275:79–109, 2002.

[5] Nelson Niu and David I. Spivak.

Polynomial functors: A mathematical theory of interaction. https://github.com/ToposInstitute/poly, October 11 2023.

[6] Jerzy Płonka.

# On a method of construction of abstract algebras.

Fundamenta Mathematicae, 61(2):183-189, 1967.

 Stanislaw Szawiel and Marek Zawadowski. Monads of regular theories. *Appl.Categor Struct*, 23:215–262, 2015.