# Compact Proofs

Measuring Quality of Understanding with a Compression-Based Metric

#### Team



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## Why care about mech interp?

- **Understanding** models
- Model evaluation and control
- Guarantees
- Model distillation

#### Why care about metrics on mech interp?

- Optimization: If we can measure it, we can optimize for it
- Automation: Suppose we got AGI tomorrow, can we build a trustworthy, automated pipeline for discovering mechanistic explanation of model behavior?
- Guarantees: Maybe mech interp can help us generate explanations with the highest standard of trustworthiness: formal proofs

## Formalizing proof length to quantify compression

Behavioral claim

True behavior of the model

theorem statement

Proof = sound computation of worst-case error (divergence in behavior)

Length of proof = cost of running computation



## Formalizing proof length to quantify compression

Behavioral claim



True behavior of the model

$$\mathbb{E}_{(x,y)\sim\mathcal{D}}\left[f(y,M(x))\right] \geq b$$

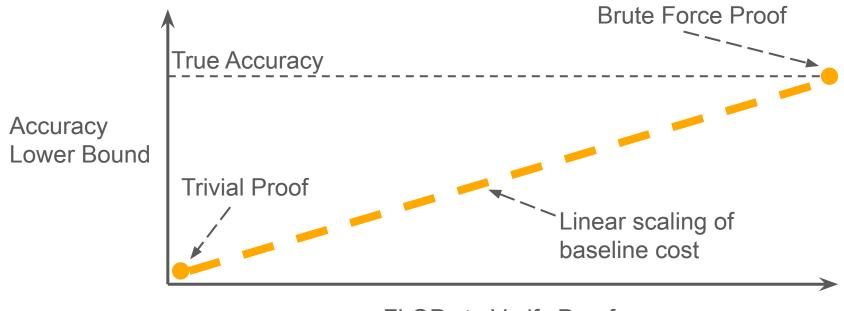
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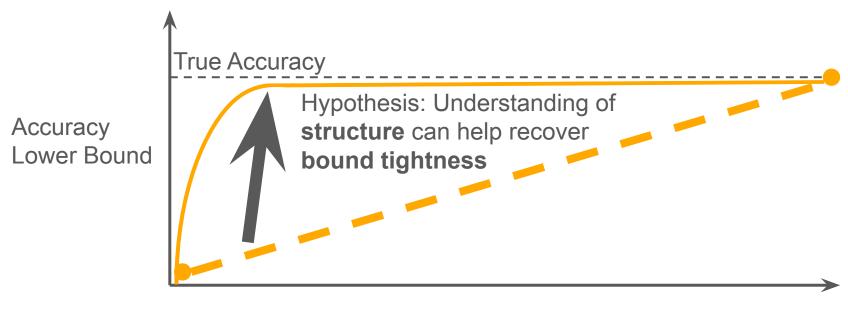


#### Quantifying the compute-cost of explanations



FLOPs to Verify Proof

#### Does understanding improve upon the linear baseline?



FLOPs to Verify Proof

## Proofs with varying mechanistic understanding

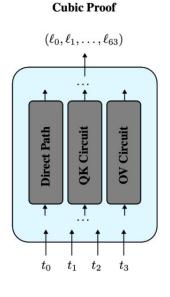
## True Model $(\ell_0,\ell_1,\ldots,\ell_{63})$ Logits Unembed **Embed** $t_1$ $t_2$ Input

FLOPs Required: Accuracy Lower Bound: Unexplained Dimension: Asymptotic Complexity:

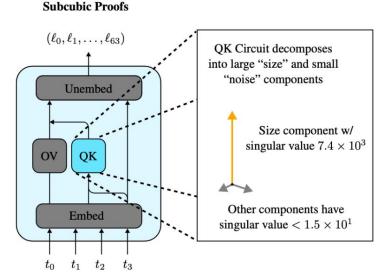


**Brute Force Proof** 

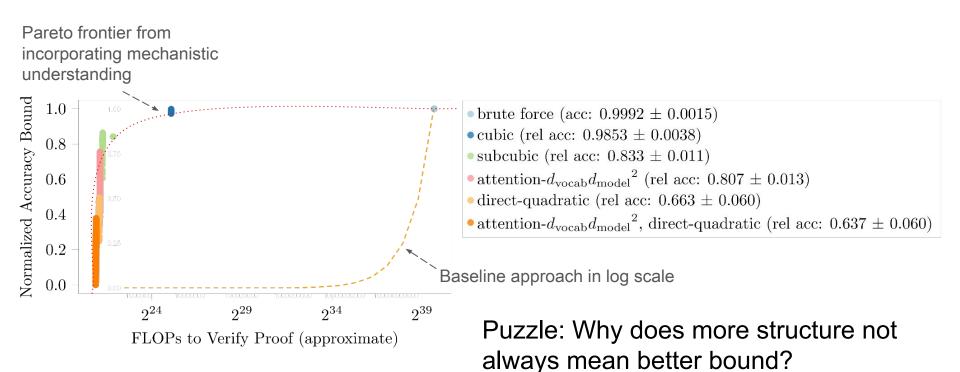
 $1.41 \times 10^{14}$  99.92%  $1.07 \times 10^{9}$   $\mathcal{O}(d_{\mathrm{vocab}}{}^{n_{\mathrm{ctx}}})$ 



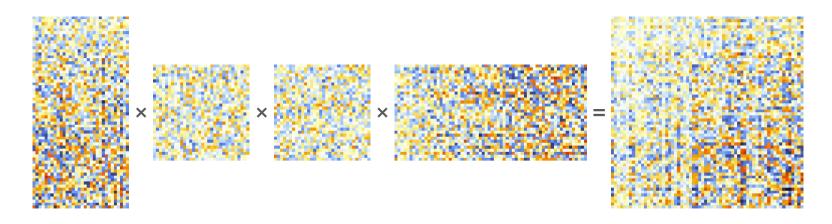




### We found an empirical "pareto frontier"



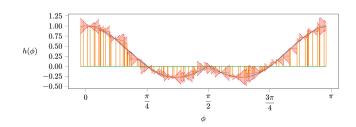
## Compounding errors from lack of structure



Approximation Strategy	Result	Complexity
(exact) max row diff	≈ 1.8	$(\mathcal{O}({d_{\mathrm{vocab}}}^2 d_{\mathrm{model}}))$
2 · (max abs value)	<b>≈</b> 2.0	$(\mathcal{O}({d_{\mathrm{vocab}}}^2 d_{\mathrm{model}}))$
max row diff on subproduct	<b>≈</b> 5.7	$(\mathcal{O}(d_{\mathrm{vocab}}d_{\mathrm{model}}^{2}))$
recursive max row diff	≈ 97	$(\mathcal{O}(d_{ ext{vocab}}d_{ ext{model}}))$

## **Applying Compact Proofs**

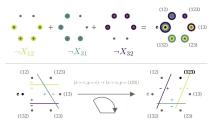
Compressing MLPs (integration)



- Ground truth for comparing mech interp approaches (groups)

Optimization targets for representation search (SAEs)





## Open Problems for Scaling Compact Proofs

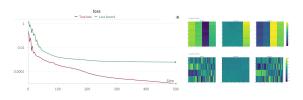
- Fix noise
  - Fine-tuning; or heuristic arguments; or sampling



- Toy model: induction heads
- Autoformalize proofs
  - AlphaProof
- Autointerp
  - Step 2: ???
- Step 3: Profit









# Questions?

## Extra Content

Overview

Proofs Approach

Case Study: Max of 4

Applications: Modular Arithmetic and SAEs

#### **Proof length**

Theorem: 
$$\mathbb{E}_{(x,y)\sim\mathcal{D}}\left[f(y,M(x))\right]\geq b$$

Our proofs consist of two components:

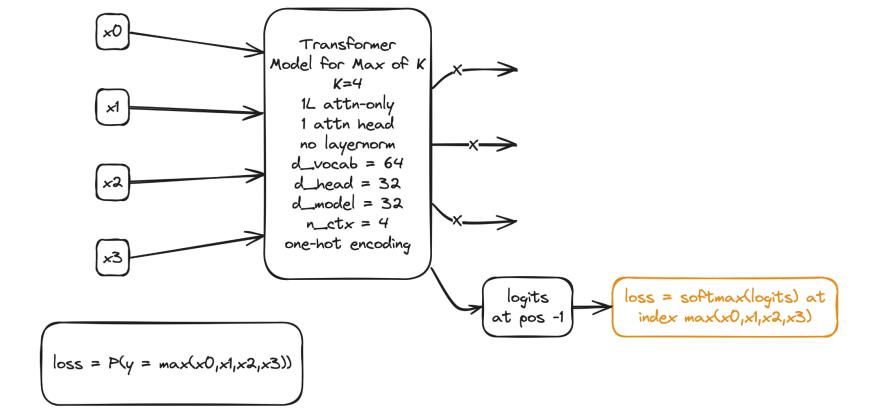
- Proof that a particular computation C, when run with any model's weights, produces a valid bound on that model's performance
- 2. A trace of running C proving that C(M) = b

In practice, proof length will be dominated by (2)

Case study: Max of K

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Proofs Approach
Case Study: Max of 4

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#### **Brute Force Proof Sketch**

```
Theorem: \mathbb{E}\left[\arg\max(M(x)[-1]) = \max_i x_i\right] > .9973

Computation C:

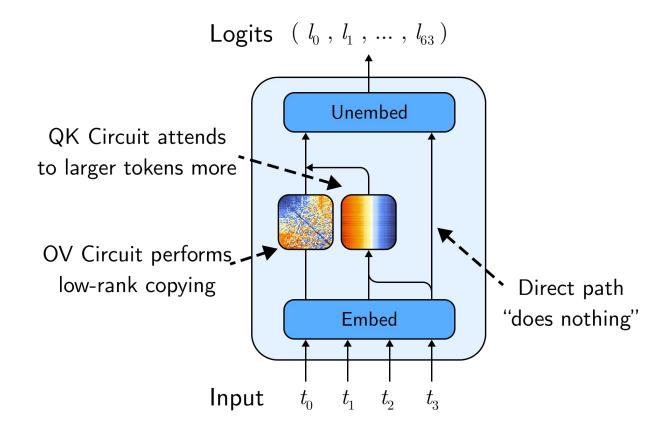
\det C(M): \\ \operatorname{count} = \emptyset \\ \operatorname{for} \times \operatorname{in} \operatorname{possible\_sequences}: \\ \operatorname{count} + = (M(x)[\ldots, -1, :].\operatorname{argmax}(\dim - 1) == x.\max(\dim - 1).\operatorname{values}).\operatorname{sum}().\operatorname{item}() \\ \operatorname{return} \operatorname{count} / \operatorname{len}(\operatorname{possible} \operatorname{sequences})
```

- 1. Lemma: *C* produces a valid bound.
  - Proof. Exercise
- 2. Lemma: C(M) = .9973

Proof: By computation.

Applications: Modular Arithmetic and SAEs

#### Attend More to Bigger Tokens & Copy



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#### Results: Proof Size vs. Tightness of Bound

Description of Proof	<b>Complexity Cost</b>	Bound
Brute force	$\mathcal{O}(v^{k+1}kd)$	$0.9992 \pm 0.0015$
Cubic	$\mathcal{O}(v^3k^2)$	$0.9531 \pm 0.0087$
Sub-cubic	$\mathcal{O}(v^2 \cdot k^2 + v^2 \cdot d)$	$0.820 \pm 0.013$
w/o mean+diff		$0.488\pm0.079$
Low-rank QK	$\mathcal{O}(v^2k^2 + vd^2 + v^2d)$	$0.795 \pm 0.014$
SVD only	QK EU&OV	$0.406\pm0.077$
Low-rank EU	$\mathcal{O}(v^2k^2 + vd + v^2d)$	$0.653 \pm 0.060$
SVD only	EU QK&OV	$(3.38\pm0.06)\times10^{-6}$
Low-rank QK&EU	$\mathcal{O}(v^2k^2 + vd^2 + vd + v^2d)$	$0.627 \pm 0.060$
SVD only	QK EU OV	$(3.38\pm0.06)\times10^{-6}$
Quadratic QK	$\mathcal{O}(v^2k^2 + \underline{vd} + \underline{v^2d})$	$0.390 \pm 0.032$
	QK EU&OV	
Quadratic QK&EU	$\mathcal{O}(v^2k^2 + \underbrace{vd}_{QK\&EU} + \underbrace{v^2d}_{OV})$	$0.285 \pm 0.036$

#### Red-Teaming Cubic Proof

#### Insight: There is a single token we want to be paying attention to (the max)

Convexity of Softmax:

(max token, non-max token)

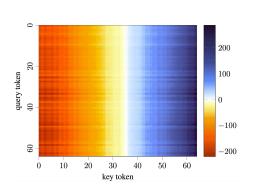
⇒ score of (how bad it is to pay attention to, how much attention is paid to it for each fixed query token)

(max token, query token)

⇒ sort non-max tokens by a combined score

If model succeeds when all non-max non-query tokens are A, it also succeeds when tokens are better than A.

Modulo handling positional encodings, this allows us to run the model on only d\_vocab<sup>3</sup> sequences instead of d\_vocab<sup>4</sup>.



#### Convexity in PyTorch Code

```
ensor,
minmax=2 attn=3 d_vocab_q d_vocab_max d_vocab_nonmax n_ctx_copies_nonmax", # noqa: F722
           The state of the crime (sin-mint one is discussed, or criticions is since-1) post-orithm attention (extrained over sequence orderings) guid to the maximum token (atth-0) and the contract of 
               Basically, this attempts to lower bound the attention that is paid to the max token and the query token, by pessinising over the order of the non-query tokens.
               Note that we could do a bit of a better job than this by taking in min_gap and a pre-computed extreme_attention matrix, if we wanted to.
               Time Complexity: O(d_vocab^3 * n_ctx^2)
                                     Topics of the streeting paid from your got buy position p

(Formal 'a 'bin [1,1], q, s, s, n come nones; ("b' from the sax token, or if the non-max

(Formal 'a 'bin [1,1], q, s, s, n come nones; ("b' from the sax token, or if the non-max

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               d_vocab, n_ctx = EQKE.shape[-1], EQKP.shape[-1]
result = torch.zeros((2, 3, d_vocab, d_vocab, d_vocab, n_ctx)).to(EQKE) + float(
                                                                       n_ctx,
       ) to[QQL)
) to[QQL)
) to[QQL)
) to[QQL)

for indices so we don't have 0 and 1 floating around

**Real **0
**Re
           for max tok in tode(range(d vocab), desc="max_tok", position=position):
    for q tok in range(max_tok + 1):
    two[:, -1] = Gut[: tok, q tok] + EQCPmi[q_tok]
    for [:, -1] = Gut[: tok, q tok] + EQCPmi[q_tok]
    for [:, -1] = Gut[: tok, q tok] + EQCPmi[q_tok]
    if q_tok = max_tok:
                                                                                                                                    continue

tmp[:,:-1] = EQKP[q_tok] + EQKE[q_tok, k_tok]

tmp.sm = (tmp / attn.scale).softmax(dim--1)

result[:, w_max, q_tok, max_tok, k_tok, 0] = tmp_sm[:,:-1].sum(
dim--1)
                                                                                                                                    )
result[:, w.nmx, q.tok, max.tok, k.tok, 0] = 0
result[:, w_qry, q_tok, max_tok, k_tok, 0] = tmp_sm[:, -1]
                                                                                                  tmp[0, :-1] = EQKP[q_tok]
tmp[1, :-1] = EQKP[q_tok].flip(dims=[0])
                                                                                                                                THE CONTROL OF THE CO
rewart; w max, q.tok, max tok, k_tok, n_copies_nommax) = tmp_set
; n=copies_max.monquery
].com(dim-1) = (tmp_set; -1) if q_tok = max tok else 0)
= comit(i, w_max, q_tok, max_tok, jok. n, copies_nommax) = result[
return result]
| - (tmp_set[:, -1] if q_tok != max_tok else 0)
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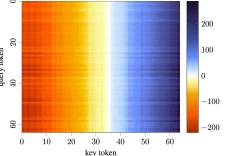
#### Red-Teaming Sub-cubic Proof

#### Insight: QK attention dominates OV

- "badness" ordering of tokens is largely independent of max token
- we can do a (slow) proof search procedure to determine "gap size" (max largest non-max)

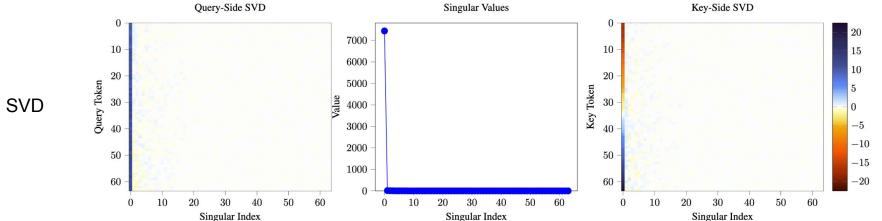
 for each query & max, we can use gap to bound how much attention we pay to non-max tokens

we then must independently pick which token is worst to parattention to (combined score is too expensive to compute exactly)



Attention Score (EQKE)

| Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (EQKE) | Solution | Score (



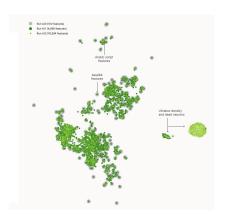
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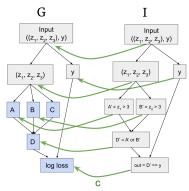
#### What we can contribute to SOTA

- Manual mech interp: where should we spend our research budget?
- SAEs: how do we reduce the dependence on human judgment?
- GPT explains GPT: we're hoping to use LLMs to formalize proofs
- Causal Scrubbing: we add rigor and a compression metric



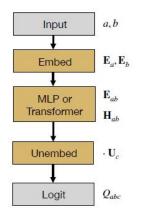






Applications: Modular Arithmetic and SAE

#### Modular Arithmetic Interpretability



**Step 1:** Embed token a and b to a circle where  $w_k = 2\pi k/p$  for some  $k \in [1, 2\cdots, p-1]$   $a \to \mathbf{E}_a \equiv (\mathbf{E}_{a,x}, \mathbf{E}_{a,y}) = (\cos(w_k a), \sin(w_k a)), b \to \mathbf{E}_b \equiv (\mathbf{E}_{b,x}, \mathbf{E}_{b,y}) = (\cos(w_k b), \sin(w_k b))$ 

#### Clock Algorithm

Step 2: compute the angle sum using multiplication.

$$\mathbf{E}_{ab} \equiv \begin{pmatrix} \mathbf{E}_{ab,x} \\ \mathbf{E}_{ab,y} \end{pmatrix} = \begin{pmatrix} \mathbf{E}_{a,x} \mathbf{E}_{b,x} - \mathbf{E}_{a,y} \mathbf{E}_{b,y} \\ \mathbf{E}_{a,x} \mathbf{E}_{b,y} + \mathbf{E}_{a,y} \mathbf{E}_{b,x} \end{pmatrix} = \begin{pmatrix} \cos(w_k(a+b)) \\ \sin(w_k(a+b)) \end{pmatrix}$$

$$\mathbf{H}_{ab} = \mathbf{E}_{ab}$$

#### Pizza Algorithm

Step 2.1: compute the vector mean.

$$\mathbf{E}_{ab} = (\mathbf{E}_a + \mathbf{E}_b)/2 = (\cos(w_k a) + \cos(w_k b), \sin(w_k a) + \sin(w_k b))/2$$

Step 2.2: using  $\mathbf{E}_{ab}$  and nonlinearities to compute  $\mathbf{H}_{ab}$ 

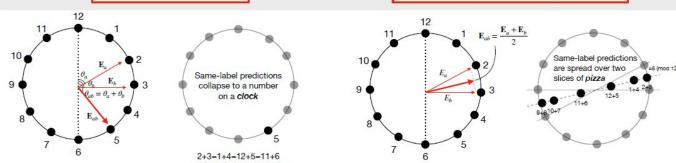
$$\mathbf{H}_{ab} = |\cos(w_k(a-b)/2)|(\cos(w_k(a+b)), \sin(w_k(a+b)))$$

**Step 3**: score possible outputs *c* using a dot product.

$$Q_{abc} = \mathbf{U}_c \cdot \mathbf{H}_{ab} , \quad \mathbf{U}_c \equiv (\mathbf{E}_{c,x}, \mathbf{E}_{c,y}) = (\cos(w_k c), \sin(w_k c))$$

 $Q_{abc}(\text{Clock}) = \cos(w_k(a+b-c))$ 





Applications: Modular Arithmetic and SAE

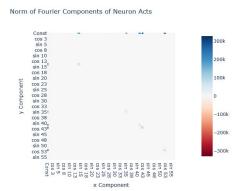
### Budget $\mathcal{O}(p^2 d_{\text{mlp}} + p^3)$

We need  $\mathcal{O}(p^2 \, d_mlp)$  just to multiply all the matrices

With this budget, we can compute everything except the MLP noise by brute force

The only interp is "unembed is low-rank"

"2D fourier basis"  $\Rightarrow$  requires  $\mathcal{O}(p^2 \, d_{\text{computation}})$  to validate



**Applications: Modular Arithmetic and SAE** 

### Explaining Pizza MLP in $\mathcal{O}(p \, d_mlp)$

#### Enough to brute force the MLP behavior (w/o noise)

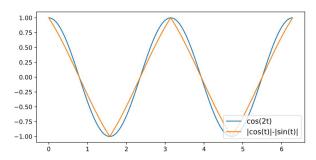


Figure 8:  $|\cos(t)| - |\sin(t)|$  is approximately  $\cos(2t)$  for any  $t \in \mathbb{R}$ 

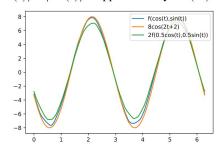


Figure 26:  $f(\cos(t), \sin(t))$  well-approximates  $8\cos(2t+2)$ .

#### L.4 Approximation

Everything becomes much clearer after realigning the matrices. For a pizza and its two corresponding principal embedding l unembedding dimensions,  $W_E'[a] + W_E'[b] \approx \cos(w_k(a-b)/2) \cdot (\cos(w_k(a+b)/2), \sin(w_k(a+b)/2))$  will be mapped by realigned  $W_1$  into its corresponding columns (which are different for every pizza), added with  $b_1$  and apply ReLU. The result will then be mapped by the realigned  $W_2$ , added with realigned  $b_2$ , and finally multipled by  $(\cos(w_kc), \sin(w_kc))$ .

For the first two principal dimensions, realigned  $W_1$  has 44 corresponding columns (with coefficients of absolute value > 0.1). Let the embedded input be  $(x,y) = W_E'[a] + W_E'[b] \approx \cos(w_k(a-b)/2) \cdot (\cos(w_k(a+b)/2) \cdot \sin(w_k(a+b)/2))$ , the intermediate columns are

 $\begin{array}{l} \text{ReLU}([0.530x-1.135y+0.253,-0.164x-1.100y+0.205,1.210x-0.370y+0.198,-0.478x-1.072y+0.215,-1.017x+0.799y+0.249,0.342x-0.048y+0.085,1.149x-0.598y+0.122,-0.443x+1.336y+0.159,-1.580x-0.000y+0.131,-1.463x+0.410y+0.178,1.038x+0.905y+0.190,0.568x+1.188y+0.128,0.235x-1.337y+0.164,-1.180x+1.052y+0.139,-0.173x+0.918y+0.148,-0.200x+1.060y+0.173,-1.342x+0.390y+0.256x+1.015x-1.246y+0.209,0.115x+1.293y+0.197,0.252x+1.247y+0.140,-0.493x+1.252y+0.213,1.120x+0.262y+0.239,0.668x+1.096y+0.205,-0.487x-1.302y+0.145,1.134x-0.862y+0.273,1.143x+0.435y+0.171,-1.285x-0.644y+0.142,-1.454x-0.285y+0.218,-0.924x+1.068y+0.145,-0.401x+0.167y+0.106,-0.411x-1.38yy+0.249,1.422x-0.117y+0.227,-0.859x-0.778y+0.121,-0.528x-0.216y+0.097,-0.884x-0.724y+0.171,1.193x+0.724y+0.131,1.086x+0.667y+0.218,0.402x+1.240y+0.213,1.069x-0.903y+0.120,0.506x-1.042y+0.153,1.404x-0.064y+0.152,0.696x-1.249y+0.199,-0.752x-0.880y+0.106,-0.956x-0.581y+0.223]). \\ 0.880y+0.106,-0.956x-0.581y+0.223]. \end{array}{}$ 

For the first principal unembedding dimension, it will be taken dot product with

 $\begin{bmatrix} 1.326, 0.179, 0.142, -0.458, 1.101, -0.083, 0.621, 1.255, -0.709, 0.123, -1.346, -0.571, 1.016, \\ 1.337, 0.732, 0.839, 0.129, 0.804, 0.377, 0.078, 1.322, -1.021, -0.799, -0.339, 1.117, -1.162, \\ -1.423, -1.157, 1.363, 0.156, -0.165, -0.451, -1.101, -0.572, -1.180, -1.386, -1.346, -0.226, \\ 1.091, 1.159, -0.524, 1.441, -0.949, -1.2481.$ 

Call this function f(x,y). When we plug in  $x=\cos(t),y=\sin(t)$ , we get a function that well-approximated  $8\cos(2t+2)$  (Figure  $\overline{26}$ ). Therefore, let  $t=w_k(a+b)/2$ , the dot product will be approximately  $8|\cos(w_k(a-b)/2)|\cos(w_k(a+b)+2)$ , or  $|\cos(w_k(a-b)/2)|\cos(w_k(a+b))$  if we ignore the phase and scaling. This completes the picture we described above.

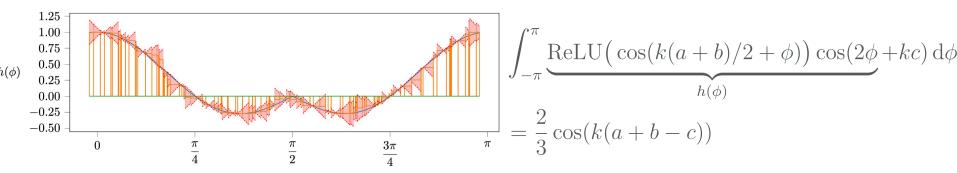
## Budget $\mathcal{O}(p + d_mlp)$ : Numerical Integration

sublinear in total parameter count ⇒ explanation w/ understanding

$$\mathcal{O}(p + d_mlp)$$
 bounds  $\Rightarrow$  evidence that numerical integration is "what's really going on"

Absolute error: 0.6 (vs 0.04 empirical error; vs 0.85 baseline)

Inability to reduce error ⇒ lacking understanding (how to align & scale intervals)



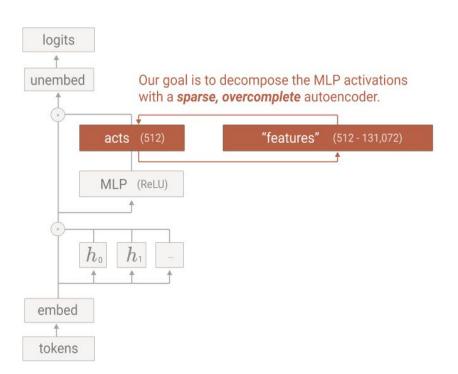
**Applications: Modular Arithmetic and SAE** 

#### Numerical Integration: Infinite-Width Limit

 $f_x(\xi)g_c(\xi)\,\mathrm{d}\xi$  $\sum_{i} f_x(\xi_i) g_c(\xi_i) w_i$  $g_c(\xi_i)w_i$  $f_x(\xi_i)$ 

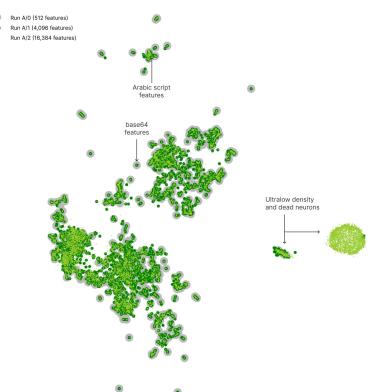
Current research question: When is this non-vacuous / interesting?

#### Sparsity Penalty in SAEs



Overview
Proofs Approach
Case Study: Max of 4

**Applications: Modular Arithmetic and SAE** 



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#### Too much human labor

#### HOW CAN WE TELL IF THE AUTOENCODER IS WORKING?

Usually in machine learning we can quite easily tell if a method is working by looking at an easily-measured quantity like the test loss. We spent quite some time searching for an equivalent metric to guide our efforts here, and unfortunately have yet to find anything satisfactory.

- - -

Thus we ended up using a combination of several additional metrics to guide our investigations:

1. Manual inspection: Do the features seem interpretable?

. . .

We think it would be very helpful if we could identify better metrics for dictionary learning solutions from sparse autoencoders trained on transformers.

#### **Tanh Penalty in Dictionary Learning**

Adam Jermyn, Adly Templeton, Joshua Batson, Trenton Bricken

. .

We found that autoencoders trained with a tanh penalty were a Pareto improvement in the space of LO and loss recovered, often by a wide margin. Unfortunately, we found that the features in these autoencoders were much harder to interpret.

https://transformer-circuits.pub/2024/feb-update/index.html#dict-learning-tanh

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#### Sparsity Penalty in SAEs

Current sparsity penalty: activations.abs().sum()

What do SAEs get you?

Case analysis structure in proofs

L1 norm: almost fine if downstream network decomposes additively not remotely useful if downstream task is non-linear

Proofs-frame suggested sparsity penalty:

```
(1 + activations.abs()).prod()
or activations.abs().log1p().sum()
```

**Applications: Modular Arithmetic and SAE** 

#### Sparsity Penalties in SAEs

