

Polynomial Functors and Shannon Entropy

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Symposium on Categorical Semantics of Entropy
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Outline

1 Introduction

- Why am I here?
- Working in the **Poly** ecosystem
- Plan of the talk

2 Background on Poly

3 Distributive functors and entropy

4 Generalizations and future work

5 Conclusion

Why am I here?

We're here to learn from each other. But what is learning?

- Somehow out of all the information out there, some of it *sticks*.
- We develop frameworks by which to *store* information.
- I'm interested in how intelligence and learning function.
- So I study how knowledge is stored and transferred in databases and...
- ...how dynamical systems interact to adapt and learn (e.g. in DNNs).

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Entropy has been put forward as an approach to intelligence and learning.

- Life can be understood as a *dissipative system*, spraying entropy.
- It does so while packing negentropy—organization—into itself.
- Polani's *empowerment* and Freer's *causal entropic forces*...
- ...are entropy-based approaches to intelligent behavior.

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I only seem to understand things when they're written categorically.

- I've been trying to understand “what entropy really is.”
- The Baez-Fritz-Leinster conception of entropy is great,...
- ...but I want to connect it in with dynamical systems or databases.

The overwhelming abundance of Poly

In January 2020 I fell in love with a category called **Poly**.

- Its applications subsume everything I'd done with categ'l databases...
- ...and everything I'd done with interacting dynamical systems.
- It's used in functional programming, type theory, higher cat'y theory.

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But it's not just very applicable, it's also very highly-structured.

- Coproducts and products that agree with usual polynomial arithmetic;
- All limits and colimits;
- At least three orthogonal factorization systems;
- A symmetric monoidal structure \otimes distributing over $+$;
- A cartesian closure q^p and monoidal closure $[p, q]$ for \otimes ;
- Another nonsymmetric monoidal structure \triangleleft that's duoidal with \otimes ;
- A left (Meyers?) \triangleleft -coclosure $\left[_ \right]$, meaning $\mathbf{Poly}(p, q \triangleleft r) \cong \mathbf{Poly}(\left[_ \right], q)$;
- An indexed right \triangleleft -coclosure, i.e. $\mathbf{Poly}(p, q \triangleleft r) \cong \sum_{f: p(1) \rightarrow q(1)} \mathbf{Poly}(p \overset{f}{\triangleleft} q, r)$;
- An indexed right \otimes -coclosure (Niu?), i.e. $\mathbf{Poly}(p, q \otimes r) \cong \sum_{f: p(1) \rightarrow q(1)} \mathbf{Poly}(p \overset{f}{\otimes} q, r)$;
- At least eight monoidal structures in total;
- \triangleleft -monoids generalize Σ -free operads;
- \triangleleft -comonoids are exactly categories; bicomodules are data migrations.

See “A reference for categorical structures on **Poly**”, arXiv: 2202.00534 [2 / 22](#)

Entropy in terms of Poly

I now use the **Poly**-ecosystem to structure my thinking.

- The abundance of structure lets me track my mental moves.
- I can check the resulting formulation using concrete examples.
- So now I try to do *everything* in **Poly**, e.g. thinking about entropy.

Entropy in terms of Poly

I now use the **Poly**-ecosystem to structure my thinking.

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- I can check the resulting formulation using concrete examples.
- So now I try to do *everything* in **Poly**, e.g. thinking about entropy.

So today I'll tell you how entropy looks from the **Poly** point of view.

- I'll show how to think of objects in **Poly** as empirical distributions.
- I'll show that there are distributive monoidal functors

$$\mathbf{Poly}^{\text{Cart}} \xrightarrow{p \mapsto \dot{p}y} \mathbf{Poly} \xrightarrow{p \mapsto (p(1), \Gamma(p))} \mathbf{Set} \times \mathbf{Set}^{\text{op}}$$

sending $p \in \mathbf{Poly}^{\text{Cart}}$ to an invariant $h(p) := (A, B) \in \mathbf{Set} \times \mathbf{Set}^{\text{op}}$.

- The Shannon entropy can then be extracted: $H(p) = \log(A^A/B)/A$.
- Properties of entropy follow from the distributive monoidality of h .

Plan

The plan for the rest of the time is as follows:

- Give background on polynomial functors.
- Explain $h: \mathbf{Poly}^{\mathbf{Cart}} \rightarrow \mathbf{Set} \times \mathbf{Set}^{\mathbf{op}}$ and its relation to entropy.
- Talk about generalizations and future work.
- Conclude.

Outline

- 1 Introduction
- 2 **Background on Poly**
 - The category **Poly**
 - Distributive monoidal structure
 - Other theoretical aspects
- 3 Distributive functors and entropy
- 4 Generalizations and future work
- 5 Conclusion

Poly for experts

What I'll call the category **Poly** has many names.

- The free completely distributive category on one object;
- The free coproduct completion of \mathbf{Set}^{op} ;
- The full subcategory of $[\mathbf{Set}, \mathbf{Set}]$ spanned by...
...functors that preserve connected limits;
- The full subcategory of $[\mathbf{Set}, \mathbf{Set}]$ spanned by coproducts of repr'bles;

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...functors that preserve connected limits;
- The full subcategory of [**Set**, **Set**] spanned by coproducts of repr'bles;
- The category of *typed sets* and colax maps between them.
 - Objects: *pairs* (S, τ) , where $S \in \mathbf{Set}$ and $\tau: S \rightarrow \mathbf{Set}$.
 - Morphisms $(S, \tau) \xrightarrow{\varphi} (S', \tau')$: *pairs* $(\varphi_1, \varphi^\#)$, where

$$\begin{array}{ccc}
 S & \xrightarrow{\varphi_1} & S' \\
 \searrow \tau & \xleftarrow{\varphi^\#} & \swarrow \tau' \\
 & \mathbf{Set} &
 \end{array}$$

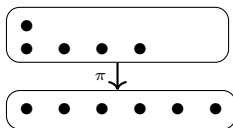
But let's make this easier.

What is a polynomial?

Algebraic

$$y^2 + 3y + 2$$

Bundle



Corolla forest

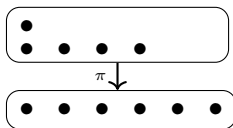


What is a polynomial?

Algebraic

$$y^2 + 3y + 2$$

Bundle



Corolla forest



You can think of the bundle as an empirical distribution:

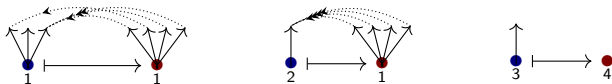
- The first outcome was drawn twice; the next three once; the rest never.
- It corresponds to the distribution $(\frac{2}{5}, \frac{1}{5}, \frac{1}{5}, \frac{1}{5}, 0, 0)$.

What is a morphism of polynomials?

Let $p := y^3 + 2y$ and $q := y^4 + y^2 + 2$



A morphism $p \xrightarrow{\varphi} q$ sends p -outcomes to q -outcomes, interpreting draws:



The category of polynomials

Easiest description: **Poly** = “sums of representable functors **Set** \rightarrow **Set**”.

- For any set S , let $y^S := \mathbf{Set}(S, -)$, the functor *represented* by S .
- Def: a polynomial is a sum $p = \sum_{i \in I} y^{p[i]}$ of representable functors.
- Def: a morphism of polynomials is a natural transformation.
- In **Poly**, usual $+$ is the coproduct and usual \times is the product.

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We will need a wide subcategory **Poly**^{Cart} \subseteq **Poly**.

- Same objects, but morphisms $p \xrightarrow{\varphi} q$ are *cartesian natural transform's*;
- ...i.e. for any function $S \rightarrow T$, the naturality square is a pullback.
- Equivalently, for each outcome $i \in p(1)$ the interpretation map

$$q[\varphi(i)] \cong p[i]$$

is a bijection. Example: there are 24 cartesian maps $y^4 \rightarrow y^4 + y^3$.

Notation

We said that a polynomial is a sum of representable functors

$$p \cong \sum_{i \in I} y^{p[i]}.$$

But note that $I \cong p(1)$. So we can write

$$p \cong \sum_{i \in p(1)} y^{p[i]}.$$

So $p(1)$ is the *outcome-set*, and elements of $p[i]$ are *draws* of outcome i .

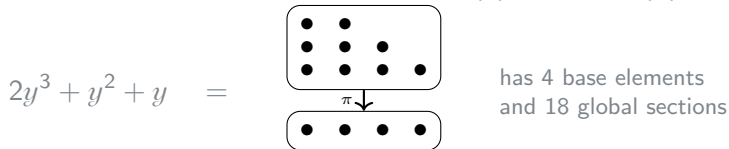
Fundamental invariants

We will be interested in two fundamental invariants of a polynomial.

- From the bundle POV, these would be *base* and *global sections*.
- So if p is represented by $E \rightarrow B$, these are B and $\mathbf{Set}_{/B}(B, E)$.
- In terms of polynomials these are

$$p(1) \cong \mathbf{Poly}(y, p) \quad \text{and} \quad \Gamma(p) := \mathbf{Poly}(p, y).$$

- E.g. for the following bundle these are $p(1) \cong 4$ and $\Gamma(p) \cong 18$.



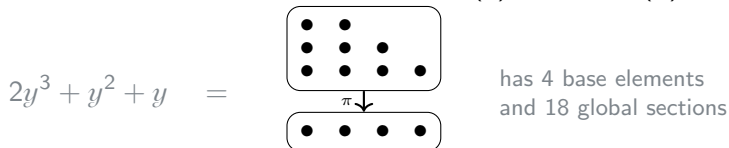
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These are functorial in opposite directions:

$$\mathbf{Poly} \xrightarrow{p \mapsto (p(1), \Gamma(p))} \mathbf{Set} \times \mathbf{Set}^{\text{op}}$$

In fact, this functor is a left adjoint, but we won't need that.

The distributive monoidal structure $(\mathbf{Poly}, 0, +, y, \otimes)$

The category **Poly** is distributive monoidal.

- The usual sum of two polynomials is their coproduct; 0 is initial.
- The usual product is the cartesian product too, but we won't use this.
- There is another operation \otimes called *Dirichlet product*. Formula:

$$p \otimes q := \sum_{(i,j) \in p(1) \times q(1)} y^{p[i] \times q[j]}$$

These are very simple bundle-wise: sum & product of base and total space:

$$\begin{array}{ccc}
 E_1 & E_2 & E_1 + E_2 \\
 \downarrow & \downarrow & \downarrow \\
 B_1 & B_2 & B_1 + B_2
 \end{array}
 +
 \begin{array}{ccc}
 E_2 & E_1 & E_1 \times E_2 \\
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- These clearly distribute: $p \otimes (q_1 + q_2) \cong (p \otimes q_1) + (p \otimes q_2)$.
- Soon we'll see how the fundamental invariants respect these oper'ns.

Derivatives and the total space

The derivative of a polynomial functor is another polynomial functor.

- Write \dot{p} for the derivative with respect to y .
- In fact, we will be much more interested in $\dot{p}y$.

$$\dot{p} = \sum_{i \in p(1)} \sum_{d \in p[i]} y^{p[i] \setminus \{d\}} \quad \text{and} \quad \dot{p}y \cong \sum_{i \in p(1)} p[i]y^{p[i]}$$

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In fact $p \mapsto \dot{p}y$ is a comonad on **Poly**^{Cart}.

- I see the counit $\dot{p}y \rightarrow p$ as “how **Poly** thinks of p as a bundle.”
- In that way p is the *base*. So let's call $p \mapsto \dot{p}y$ the *total space functor*.

Bifibration $\mathbf{Poly} \rightarrow \mathbf{Set}$

The last theory we'll need is the bifibration $\mathbf{Poly} \rightarrow \mathbf{Set}$.

- The functor $p \mapsto p(1)$ is both a fibration and an op-fibration.
- In fact it's even more: a distributive monoidal $*$ -bifibration!

Down to earth what does this mean? Let p be a polynomial and A a set.

- For any function $f: A \rightarrow p(1)$ we can take the pullback

$$\begin{array}{ccc}
 f^*p & \longrightarrow & p \\
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- So f^* takes polys with outcome-set $p(1)$ to those with outcome-set A .
- This operation has both a left adjoint $f_!$ and a right adjoint f_* .

$$f_!p := \sum_{b \in B} y \prod_{a \rightarrow b} p[a] \quad \text{and} \quad f_*p := \sum_{b \in B} y \sum_{a \rightarrow b} p[a]$$

- I.e., for any $f: A \rightarrow B$, we have $\mathbf{Poly}_A(f^*q, p) \cong \mathbf{Poly}_B(q, f_*p)$.

As we lump outcomes together, we add up the draws.

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2 Background on Poly

3 **Distributive functors and entropy**

- Distributive functor $\mathbf{Poly}^{\mathbf{Cart}} \rightarrow \mathbf{Set} \times \mathbf{Set}^{\text{op}}$
- Entropy and entropy density

4 Generalizations and future work

5 Conclusion

Total space as distributive

We're now ready to get to work on how all this relates to entropy.

- The approach is to extract two invariant sets from any polynomial.
- This process is “good” in that it is a distributive monoidal functor.
- We'll extract the extensive and intensive entropies from these.

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So first, we want to see that $p \mapsto \dot{p}y$ is distributive monoidal.

- The derivative is linear, $(p + q) \dot{=} \dot{p} + \dot{q}$, and so is $p \mapsto py$.
- So $p \mapsto \dot{p}y$ preserves coproducts. What about \otimes ?

$$\begin{aligned}
 (\dot{p}y) \otimes (\dot{q}y) &\cong \sum_{i \in p(1)} p[i]y^{p[i]} \otimes \sum_{j \in q(1)} q[j]y^{q[j]} \\
 &\cong \sum_{(i,j) \in p(1) \times q(1)} p[i] \times q[j]y^{p[i] \times q[j]} \\
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So $(\mathbf{Poly}^{\text{Cart}}, 0, +, y, \otimes) \xrightarrow{p \mapsto \dot{p}y} (\mathbf{Poly}, 0, +, y, \otimes)$ is distributive monoidal.

Fundamental invariants as distributive

The fundamental invariants $p \mapsto (p(1), \Gamma(p))$ are also distributive

$$(\mathbf{Poly}, 0, +, y, \otimes) \xrightarrow{p \mapsto (p(1), \Gamma(p))} (\mathbf{Set} \times \mathbf{Set}^{\text{op}}, (0, 1), +, (1, 1), \otimes)$$

But what exactly is all this structure on $\mathbf{Set} \times \mathbf{Set}^{\text{op}}$?

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- $\mathbf{Set} \times \mathbf{Set}^{\text{op}}$ has coproducts: $(A_1, B_1) + (A_2, B_2) \cong (A_1 + A_2, B_1 \times B_2)$.
- It has another symmetric monoidal structure with unit $(1, 1)$:

$$(A_1, B_1) \otimes (A_2, B_2) := (A_1 \times A_2, B_1^{A_2} \times B_2^{A_1})$$

- And these distribute “because” $B^{A_1 + A_2} (B_1 B_2)^A \cong (B^{A_1} B_1^A) (B^{A_2} B_2^A)$.

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- And these distribute “because” $B^{A_1+A_2}(B_1 B_2)^A \cong (B^{A_1} B_1^A)(B^{A_2} B_2^A)$.

Why do the fundamental invariants (as a pair) preserve $+$ and \otimes ?

- We have $(p + q)(1) \cong p(1) + q(1)$ and $\Gamma(p + q) \cong \Gamma(p) \times \Gamma(q)$.
- This says they preserve $+$. One also checks they preserve \otimes :

$$(p \otimes q)(1) \cong p(1) \times q(1) \quad \text{and} \quad \Gamma(p \otimes q) \cong \Gamma(p)^{q(1)} \times \Gamma(q)^{p(1)}$$

Taking stock

Let's denote the composite of our distributive functors by \hat{h} :

$$\begin{array}{ccc}
 \mathbf{Poly}^{\mathbf{Cart}} & \xrightarrow{p \mapsto \dot{p}y} & \mathbf{Poly} & \xrightarrow{p \mapsto (p(1), \Gamma(p))} & \mathbf{Set} \times \mathbf{Set}^{\mathbf{OP}} \\
 & & \parallel & & \nearrow \\
 & & \hat{h} & &
 \end{array}$$

- The claim is that \hat{h} extracts everything you need to calculate entropy.
- Preserving $+$ and \otimes gives us properties of entropy.

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$$\underbrace{\hspace{15em}}_{\hat{h}}$$

- The claim is that \hat{h} extracts everything you need to calculate entropy.
- Preserving $+$ and \otimes gives us properties of entropy.

Define a real number $L(A, B) := \frac{\log(\#A^{\#A}) - \log(\#B)}{\#A}$. Then:

Theorem

Let p be a polynomial, considered as a probability distribution P , and let $H(P)$ be its Shannon entropy. Then we have

$$H(P) = L(\hat{h}(p))$$

The categorical partition function and entropy

I'm unfamiliar with the thermo picture. Joint with James Dama:

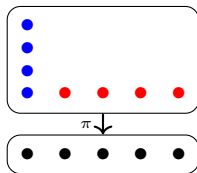
- One should think of Shannon entropy as an *entropy density*.
- The thermo picture defines a *partition function* Ω for distributions.
- For $p \in \mathbf{Poly}$ with $h(p) = (A, B)$, this would be $\Omega_p := \frac{A^A}{B}$.
- Then the extensive entropy of p is given by $E(p) := \log \Omega(p)$.
- And the Shannon entropy of p is the density $H(p) := E(p)/A$.

The categorical partition function and entropy

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- And the Shannon entropy of p is the density $H(p) := E(p)/A$.

For example consider the bundle for $p := y^4 + 4y^1$:



- We find $\dot{p}y = 4y^4 + 4y$, so $h(p) = (4 + 4, 4^4) = (8, 4^4)$.
- So $\Omega_p = \frac{8^8}{4^4}$, Ext've: $E(p) = \log \Omega_p = 16$, Shannon: $H(p) = 16/8 = 2$.

Consequences of distributivity

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- Suppose you write p as a sum, $p := \sum_{a \in A} p_a$.
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It follows from the fact that h preserves sums that:

$$\Omega_p = \Omega_{f_*p} \times \prod_{a \in A} \Omega_{p_a} \quad \text{and} \quad E(p) = E(f_*p) + \sum_{a \in A} E(p_a)$$

The usual “chain rule” for Shannon entropy follows directly from this.

A geometric viewpoint on Shannon entropy

Think of objects $(A, B) \in \mathbf{Set} \times \mathbf{Set}^{\text{op}}$ as representing *formal rectangles*.

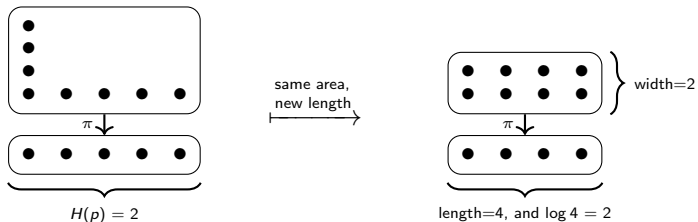
- Here A is its area, $\sqrt[A]{B}$ is its width, and $A/\sqrt[A]{B}$ is its length.
- Adding two rectangles $(A_1, B_1) + (A_2, B_2) = (A_1 + A_2, B_1 \times B_2)$...
- ...add the areas and take the weighted geometric mean of the widths.
- Multiplying two rectangles $(A_1, B_1) \otimes (A_2, B_2) = (A_1 A_2, B_1^{A_2} B_2^{A_1})$...
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- The log-length $\log(A/\sqrt[A]{B})$ of the rectangle is the Shannon entropy.

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Let $p := y^4 + 4y$, so $h(p) = (8, 4^4)$. Width = $\sqrt[8]{4^4} = 2$, length = $8/2 = 4$.

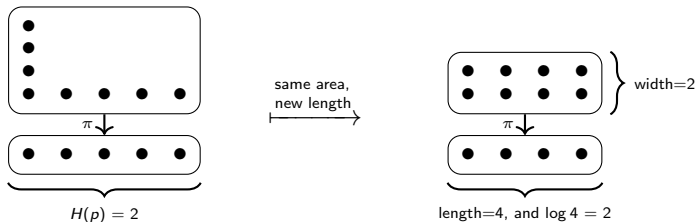


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If $q = Ly^W$ was rectangular to begin with, it'll stay that way.

Outline

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- 2 Background on Poly
- 3 Distributive functors and entropy
- 4 Generalizations and future work**
 - Functoriality?
- 5 Conclusion

Functoriality?

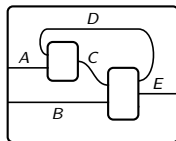
One of the big questions for me is: what does functoriality buy you?

- The distributivity of $h: \mathbf{Poly}^{\mathbf{Cart}} \rightarrow \mathbf{Set} \times \mathbf{Set}^{\mathbf{op}}$ means something.
- It gives us well-known facts about entropy and entropy density.
- But what about the fact that h is functorial?
- Logarithms have no clue about what maps in $\mathbf{Set} \times \mathbf{Set}^{\mathbf{op}}$ mean.

Entropy and dynamics?

Returning to my goals, I'd like to understand learning.

- If entropy will be involved, I want it to be about dynamical systems.
- The groupoid $\dot{p}y$ is kind of dynamic: little $p[i]$'s spinning around.
- But what about the *point* of Shannon entropy: communication?



- Shouldn't we be able to see Huffman coding or something here...?
- What about "empowerment" or "causal entropic forces"?

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- Entropy is well-known throughout the scientific and technical world.
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- Both have applications to interacting dynamical systems.

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There is (at least one) interpretation of entropy within **Poly**.

- Objects in **Poly** can be viewed as empirical distributions.
- There is a $(+, \otimes)$ -preserving functor $h: \mathbf{Poly}^{\text{Cart}} \rightarrow \mathbf{Set} \times \mathbf{Set}^{\text{op}}$.
- If $h(p) = (A, B)$ then $H(p) = \log(A / \sqrt[A]{B}) / A$.
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- So all the entropy-relevant data of p is encapsulated in two sets.

Entropy still feels somehow foreign to me.

- Hopefully, having different categorifications will help clarify it.
- I still have hope that it will bond with the dynamics of **Poly**.

Thanks! Comments and questions welcome...